

# Building Cyclic Data in a Functional-Like Language Extended with Monotonic Objects

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# Part of General Problem

## Building Cyclic Data in a Functional-Like Language Extended with Monotonic Objects

is a particular task within the following general problem

Construction of an **intermediate** model of computation **between** the **functional** and **object-oriented** programming paradigms, **preserving the majority of nice properties of the functional one** and **extending the domain** of its efficient applications

What is **the main limitation of functional** programming?

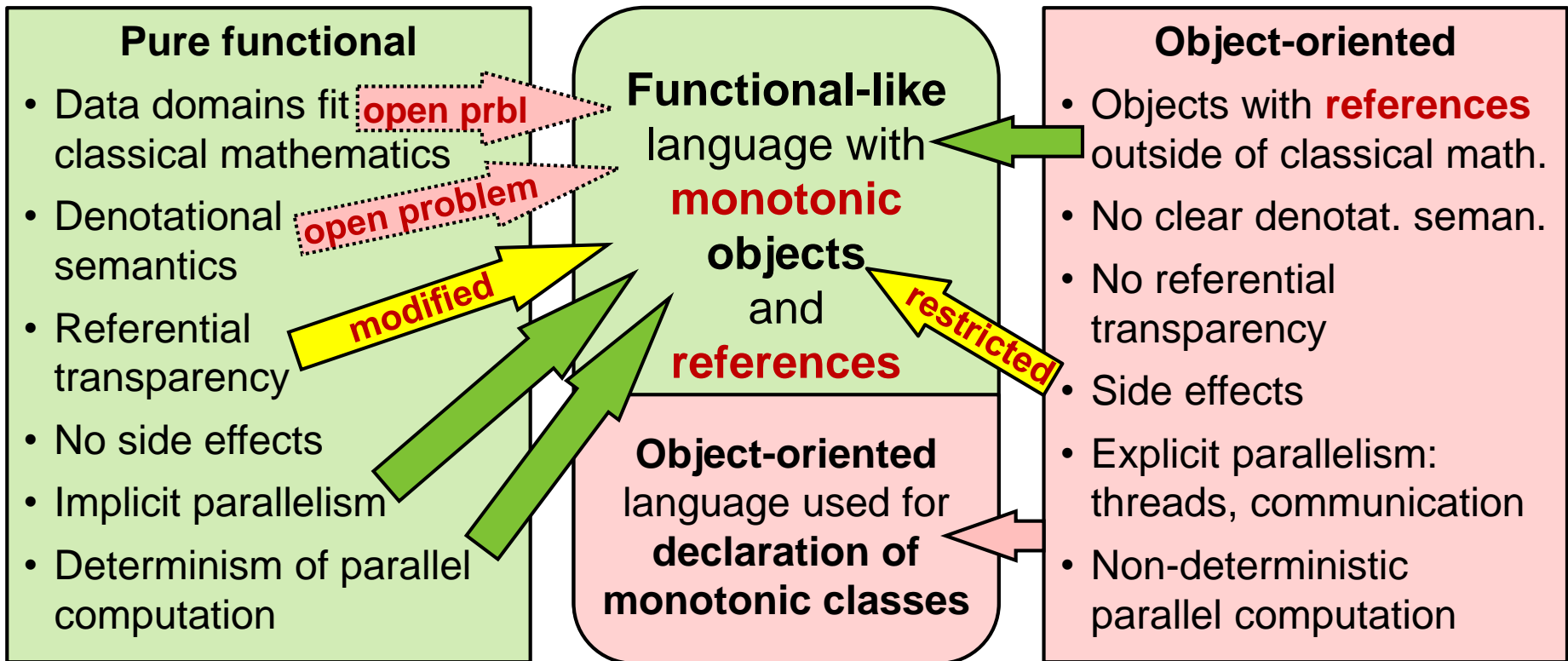
- **Functional** languages allow us to declare and efficiently manipulate **tree data only**, by composing terms of constructors
- **Object-oriented** languages allow us to efficiently represent and traverse **arbitrary graphs**, denoting relations between vertices and/or vertices and edges by **references** to objects

Hence the task

# General Problem Statement

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is a particular task within the following general problem



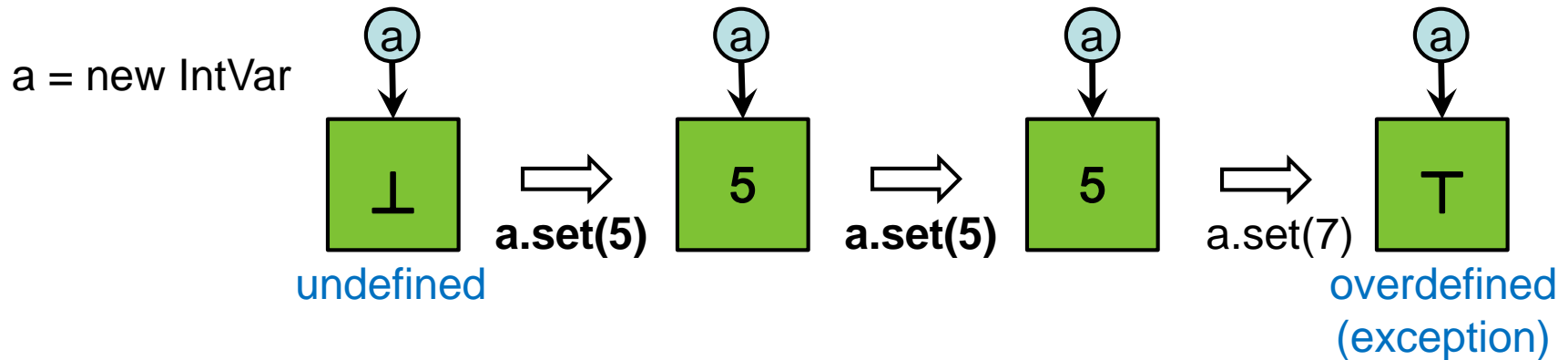
# Monotonic Objects and Classes

**Definition.** **Monotonic classes and objects** are such that **functional-like programs invoking methods on them** satisfy the following properties

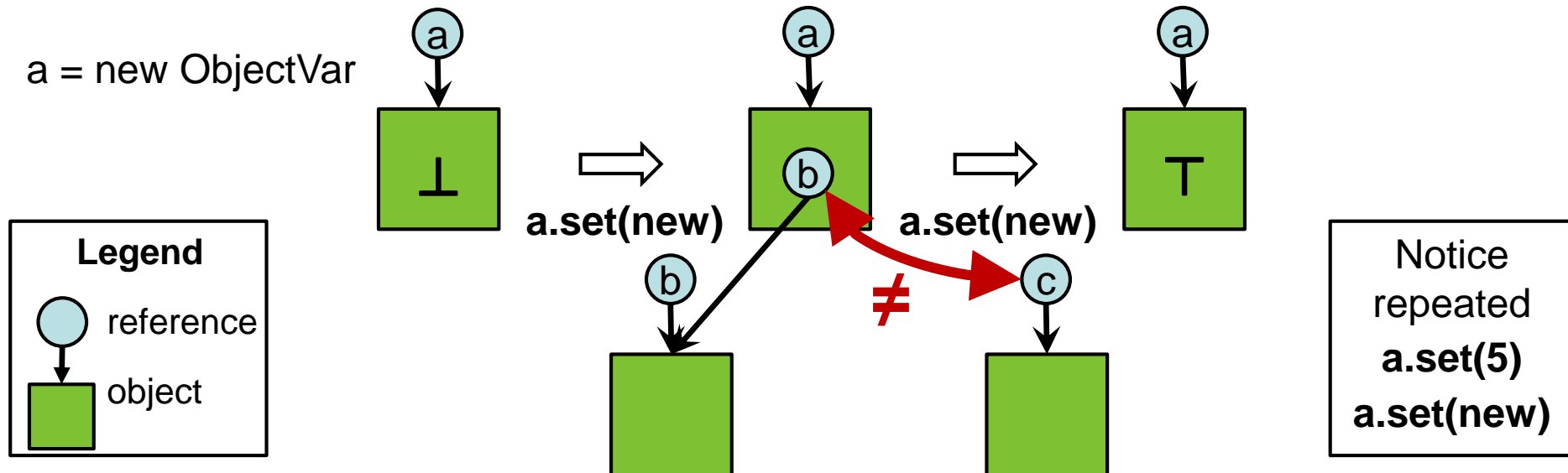
- **Operational properties** (formal)
  - **Determinism** of parallel computation:
    - results obtained in different order of computation are **equivalent** (their difference is unobservable)
  - **Idempotency**:
    - recomputation of an expression produces an **equivalent result** and **side effect difference is unobservable**
- **Target properties** (still informal)
  - Existence of **semilattices** in which objects change **monotonically**
    - derived from declarations of monotonic classes
  - Existence of **denotational semantics**
    - generalizing that of pure functional languages
    - not using the order of computation (parallel, in essence)

# Simplest Monotonic Class: Arvind's I-Structure

## Monotonic with a primitive value



## Non-monotonic with a reference value



# Simplest Monotonic Class: Java Code

```
public class IntVar {  
  
    boolean defined = false;  
    int value;  
  
    public synchronized int get() {  
        if (!defined) wait();  
        return value;  
    }  
  
    public synchronized void set(int x) {  
        if (!defined) {  
            value = x;  
            defined = true;  
            notifyAll();  
        }  
        else if (value != x)  
            throw new RuntimeException();  
    }  
}
```

**Monotonic** class with a value of  
primitive type int

```
public class ObjectVar {  
  
    boolean defined = false;  
    Object value;  
  
    public synchronized Object get() {  
        if (!defined) wait();  
        return value;  
    }  
  
    public synchronized void set(Object x) {  
        if (!defined) {  
            value = x;  
            defined = true;  
            notifyAll();  
        }  
        else if (value != x)  
            throw new RuntimeException();  
    }  
}
```

**Non-monotonic** class with a value of  
reference type Object

# Monotonic I-Structure with a Reference Value

## Monotonic with a reference value

① `a = new ObjectVar`

a

d

④ `d = a.get`

A new intermediate object is created on the first `set`

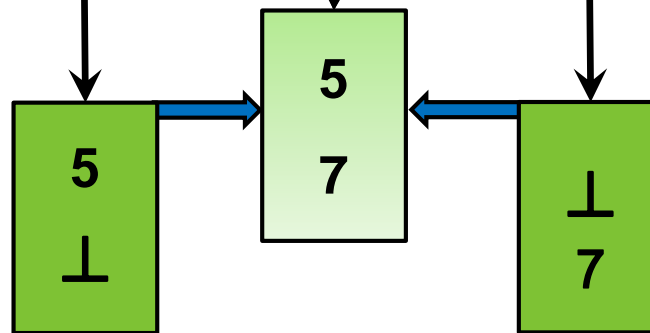
The intermediate object is returned by `get`

② `b = new ...  
a.set(b)`

b

③ `c = new ...  
a.set(c)`

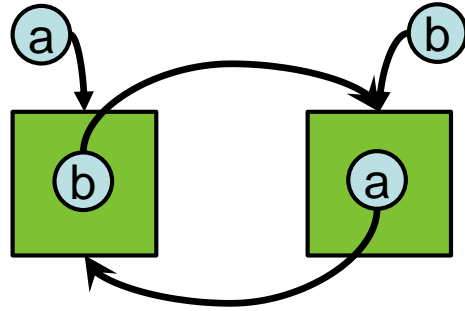
c



**Unification of the set objects**

# How to Monotonically Create a Cycle (1)

**Our wish**

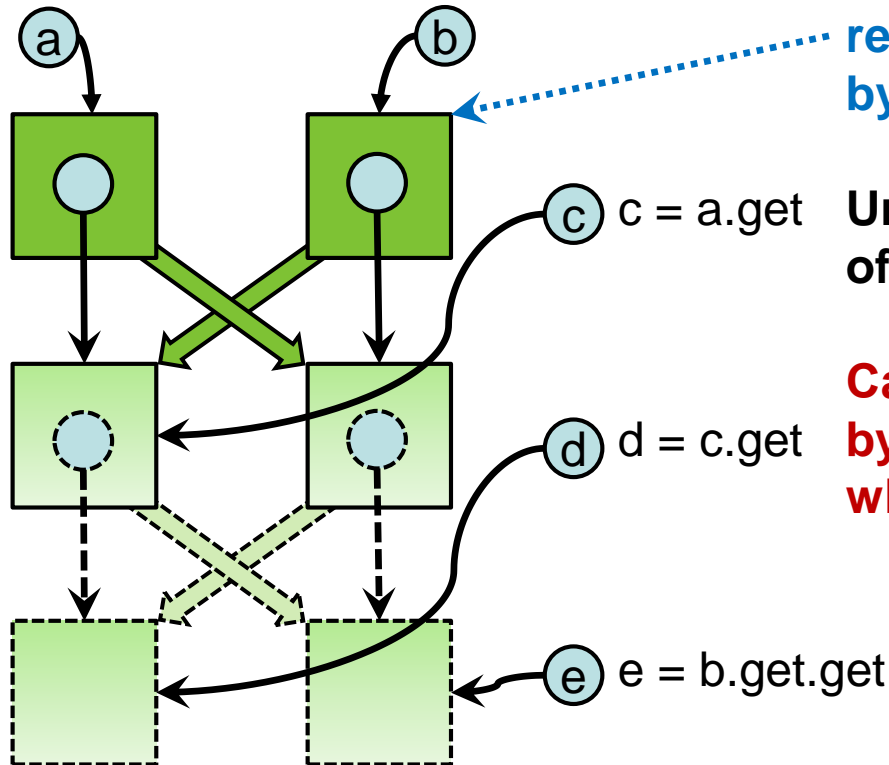


```
a = new  
b = new  
a.set(b)  
b.set(a)
```

Easy in an object-oriented language  
**but it is non-monotonic**

**Monotonic  
1<sup>st</sup> solution**

(Example 4)



Solid-line objects and references are created by the above code

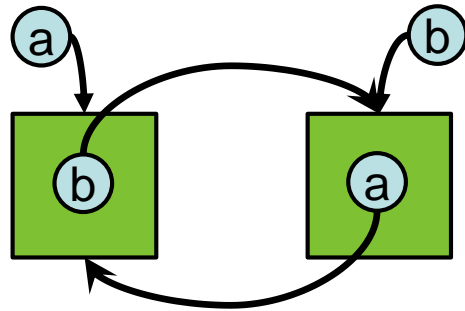
Undesirable unrolling of the cyclic graph

**Can't catch a cycle by comparing references while traversing**



# How to Monotonically Create a Cycle (2)

**Our wish**

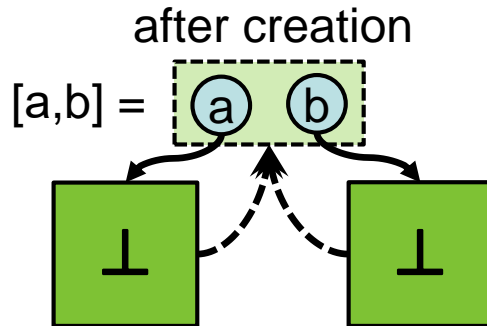


```
a = new  
b = new  
a.set(b)  
b.set(a)
```

Easy in an object-oriented language  
**but it is non-monotonic**

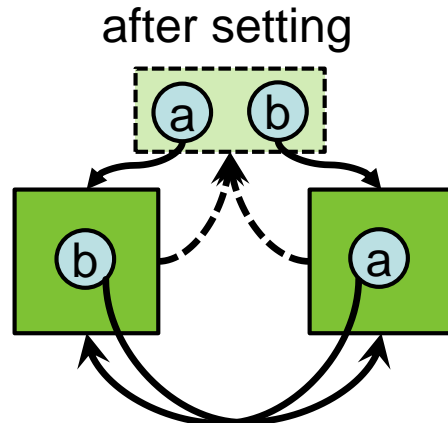
**Monotonic  
2<sup>nd</sup> solution**

(Example 5a)



```
[a,b] = new[2]
```

Simultaneous creation  
of objects a and b



```
a.set(b)  
b.set(b)
```

Objects know references  
to each other from birth

**In this case  
plain setting as in OOL  
does not violate  
monotonicity**

# How to Monotonically Create a Cycle (3)

We know **one more solution**\* to the **problem of building a cyclic data structure**, which is simpler to express as an algorithm than to draw a picture:

## Algorithm

1. Consider building a cycling graph **as in solution 1**
2. Let us have **one object as the root**
3. Invoke on the root the **minimize** method that waits for all fields of the objects accessible from the root to **become defined** and then **merges the objects that are indistinguishable** by all operations
4. The operation **minimize** returns a **fresh reference to the root** such that all **objects accessible** from it get a fresh **unique reference** for each

**Monotonic  
3<sup>rd</sup> solution**

(Example 5b)

**In the 2<sup>nd</sup> and 3<sup>rd</sup> solutions catching cycles is possible like in a plain object-oriented language**

\*There may be other solutions which we don't know yet

# Conclusion

- We have made first steps to construction of a **new model of computation** intermediate **between functional and object-oriented paradigms**
- A programming language that implements the model is **two-level**:
  - The **higher** level is a **functional-like** language
  - The **lower** level is a common **object-oriented** language
- The **main idea** is to **restrict the methods** of the classes (called **monotonic**) that are **used in the functional language** in such a way that the following **properties of functional programs** are **preserved**:
  - **Determinism** of parallel computation
  - **Idempotency** of side effects and results
- The key initial problem to be solved is **overcoming the main limitation of functional languages** that **only tree-like structures** can be constructed
- We have demonstrated **examples of monotonic classes**, by using which a **functional program can build cyclic structures** where **each object has a programmatically accessible unique reference**